Simultaneous recovery of the consensus and structure of permutations

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UW Theory Seminar 2/2/16

The "Sushi preference" data

N = 5000 people ranked n = 12 types of sushi

sake |ebi |ika |uni |tamago |kappa-maki |tekka-maki |anago |toro |maguro ebi |kappa-maki |tamago |ika |toro |maguro |tekka-maki |anago |sake |uni toro |ebi |maguro |ika |tekka-maki |uni |sake |anago |kappa-maki |tamago tekka-maki |tamago |sake |ebi |ika |kappa-maki |maguro |toro |uni |anago tamago |maguro |kappa-maki |ebi |sake |anago |uni |tekka-maki |toro |ika uni |toro |ebi |anago |maguro |tekka-maki |ika |sake |kappa-maki |tamago maguro |ika |toro |tekka-maki |ebi |uni |sake |tamago |anago |kappa-maki

Consensus Ranking Problem

Given a set of rankings $\{\pi_1,\pi_2,\dots\pi_N\}\subset\mathbb{S}_n$ find the consensus ranking π_0 such that

$$\pi_0 = \underset{\mathbb{S}_n}{\operatorname{argmin}} \sum_{i=1}^N d(\pi_i, \pi_0)$$

for d= inversion distance / Kendall au-distance / "bubble sort" distance

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Consensus Ranking Problem

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This problem is NP-hard []

Related work

Consensus Ranking/Single parameter/Mallows model

[Cohen, S, Singer 99] CSS ALGORITHM = greedy search on Q improved by extracting strongly connected components

[Ailon,Newman,Charikar 05] Randomized algorithm guaranteed 11/7 factor approximation (ANC)

[Mohri, Ailon 08] linear program

[Mathieu, Schudy 07] $(1+\epsilon)$ approximation, time $\mathcal{O}(n^6/\epsilon+2^{2^{O(1/\epsilon)}})$

[Davenport, Kalagnanan 03] Heuristics based on edge-disjoint cycles used by our B&B implementation

[Conitzer,D,K 05] Exact algorithm based on integer programming, better bounds for edge disjoint cycles (DK)

[Betzler,Brandt, 10] Exact problem reductions [Awasthi,Blum,Sheffet,Vijayaraghavan 14]

Most of this work based on the MinFAS view

$$Q_{ij} > .5 \Leftrightarrow i \bullet \stackrel{Q_{ij} - .5}{\longrightarrow} \bullet j$$

Prune graph to a DAG removing minimum weight

Extensions and applications to social choice

Social choice

- ► Inferring rakings under partial and aggregated information [ShahJabatula08], [JabatulaFariasShah10]
- ▶ Vote elicitation under probabilistic models of choice [LuBoutillier11]
- Voting rules viewed as Maximum Likelihood [ConitzerSandholm08]
- Algorithms guaranteed to retrive certain "winners" [LinAgarwal14]
 "Noisy sorting"
- ▶ Using Hodge decompositions and L1, L2 distances [JiangLimYaoYe11]
- ► Noisy comparison [BravermanMossel08]

$ML\ Estimation/Multiple\ parameters/GM\ model$

I[VlignerVerducci 86] $\vec{\theta}$ estimation; heuristic for π_0

FV ALGORITHM/BORDA RULE

- 1. Compute $ar{q}_j, j=1:n$ column sums of Q
- 2. Sort $(\bar{q}_j)_{j=1}^n$ in increasing order; π_0 is sorting permutation
- $ightharpoonup ar{q}_i$ are Borda counts
- ► FV is consistent for infinite *N*



Generalizing consensus ranking

▶ Not all inversions are equally important

Sushi preferences for uni have no consensus

sake |ebi |ika |uni |tamago |kappa-maki |tekka-maki |anago |toro |maguro ebi |kappa-maki |tamago |ika |toro |maguro |tekka-maki |anago |sake |uni toro |ebi |maguro |ika |tekka-maki |uni |sake |anago |kappa-maki |tamago |tekka-maki |tamago |sake |ebi |ika |kappa-maki |maguro |toro |uni |anago |tamago |maguro |kappa-maki |ebi |sake |anago |uni |tekka-maki |toro |ika |uni |toro |ebi |anago |maguro |tekka-maki |ika |sake |kappa-maki |tamago |maguro |tekka-maki |toro |tekka-maki |ebi |uni |sake |tamago |anago |kappa-maki |

Generalizing consensus ranking

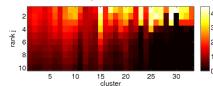
▶ Not all inversions are equally important

... but there is consensus for maguro (tuna) and tekka-maki (tuna roll) sake |ebi |ika |uni |tamago |kappa-maki |tekka-maki |anago |toro |maguro ebi |kappa-maki |tamago |ika |toro |maguro |tekka-maki |anago |sake |uni toro |ebi |maguro |ika |tekka-maki |uni |sake |anago |kappa-maki |tamago tekka-maki |tamago |sake |ebi |ika |kappa-maki |maguro |toro |uni |anago tamago |maguro |kappa-maki |ebi |sake |anago |uni |tekka-maki |toro |ika uni |toro |ebi |anago |maguro |tekka-maki |ika |sake |kappa-maki |tamago maguro |ika |toro |tekka-maki |ebi |uni |sake |tamago |anago |kappa-maki

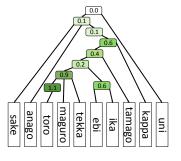
Generalizing consensus ranking

 \blacktriangleright Not all inversions are equally important introduce importance/weight parameters $\vec{\theta}$

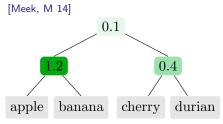
Irish College Admissions data
Parameters of top 10 ranks in the 33 largest clusters found



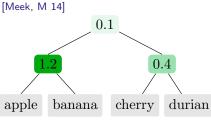
Combinatorial structure present



described by a tree



```
	au= tree structure \pi_0(	au)= induced central ranking 	heta_{1:n-1}= parameters at nodes
```



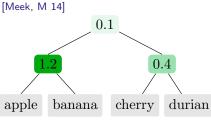
au= tree structure

 $\frac{\pi_0(\tau)}{2}$ = induced central ranking

 $\theta_{1:n-1} = \text{parameters at nodes}$ Inversions are penalized by θ_i parameters

Example: $\vec{ heta} = (0.1, 1.2, 0.4)$

$$Cost(a|b|c|d) = 0$$



au = tree structure

 $\pi_0(au) = ext{induced central ranking} \ heta_{1:n-1} = ext{parameters at nodes}$

Inversions are penalized by θ_i parameters

Example: $\vec{\theta} = (0.1, 1.2, 0.4)$

Cost(a|b|c|d) = 0

Cost(b|a|c|d) = 1.2

[Meek, M 14]

0.1

apple banana cherry durian

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 $Cost(c|b|a|d) = 1.2 + 2 \times 0.1$

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 $Cost(c|b|a|d) = 1.2 + 2 \times 0.1$

$$P(a|b|c|d) \propto e^{0}$$

 $P(b|a|c|d) \propto e^{-1.2}$
 $P(c|b|a|d) \propto e^{-1.2-2\times0.1}$

RIM distribution $P_{\tau, \vec{\theta}}$

Let $v_i =$ number of inversions of π at node i

$$P_{\boldsymbol{\tau},\vec{\theta}}(\pi) \propto \prod_{i \in nodes} \exp(-\theta_i \mathbf{v}_i)$$

[Meek, M 14] 0.10.4apple banana cherry durian

 $\pi_0(\tau) = \text{induced central ranking}$ $\theta_{1:n-1} = \mathsf{parameters} \ \mathsf{at} \ \mathsf{nodes}$ Inversions are penalized by θ_i parameters

Example:
$$\vec{\theta} = (0.1, 1.2, 0.4)$$

$$\mathsf{Cost}(a|b|c|d) = 0$$

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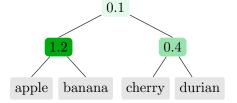
Normalization constant

$$Z(\tau, \theta) = \prod_{i \in nodes} G(L_i, R_i, \exp(-\theta_i))$$

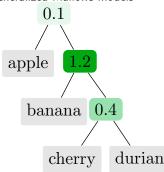
with
$$G(L, R, q) = \frac{(q)_{L+R}}{(q)_L(q)_L}$$
, $(q)_n = \prod_{i=1}^n (1 - q^i)$.

Structure τ known as Riffle Independence model [Huang, Guestrin 12]

The RIM is a general flexible model



- ▶ any tree structure
- ▶ any parameters (but $\theta_i \ge 0$ suffices)
- ▶ includes the Mallows and Generalized Mallows models



Max Likelihood Estimation for RIM

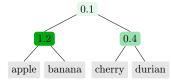
[M,Meek 14]

▶ Problem Given permutations $\pi_1, \dots \pi_N$, infer τ, θ

Max Likelihood Estimation for RIM

[M, Meek 14]

- ▶ Problem Given permutations $\pi_1, \dots \pi_N$, infer τ, θ
- ▶ Identifiability of θ
 - ▶ reorder to obtain cannonical representation, with $\theta_i \ge 0$ for all $i \in nodes$
 - ightharpoonup given au, $heta_i$ can be estimated by convex univariate minimization

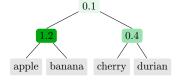


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Max Likelihood Estimation for RIM

[M, Meek 14]

- ▶ Problem Given permutations $\pi_1, \ldots \pi_N$, infer τ, θ
- ▶ Identifiability of θ
 - reorder to obtain cannonical representation, with $\theta_i > 0$ for all $i \in nodes$
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Identifiability of τ

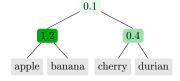
Theorem[M, Meek 14] A model τ , θ is identifiable iff

- 1. $\theta_i > 0$ for all $i \in nodes$
- 2. $\theta_i \neq \theta_{pa(i)}$ for all $i \in nodes$ (pa(i)) is the parent of node i in τ)

Max Likelihood Estimation for RIM

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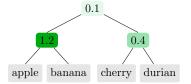


Identifiability of au

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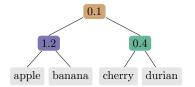
- 1. $\theta_i > 0$ for all $i \in nodes$
- 2. $\theta_i \neq \theta_{pa(i)}$ for all $i \in nodes$ (pa(i) is the parent of node i in τ)
- \blacktriangleright Hardness of τ estimation
 - Estimating π_0 is NP-hard [Duchi, Mackey, Jordan 13]
 - Estimating τ structure given π_0 is tractable

Sufficient statistics



Q(d a b c) =	a	Ь	С	d	
	_	1	0	0	a
	0	_	1	0	Ь
	0	0	_	0	С
	1	1	1	_	d

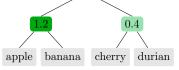
Sufficient statistics



$$Q(d|a|b|c) = \begin{bmatrix} a & b & c & d \\ - & 1 & 1 & 0 & a \\ 0 & - & 1 & 0 & b \\ \hline 0 & 0 & - & 0 & c \\ 1 & 1 & 1 & - & d \end{bmatrix}$$

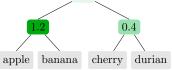
$$Cost(d|a|b|c) = 0.1 \times 2 + 1.2 \times 0 + 0.4 \times 1$$

$\underset{0.1}{\text{Max Likelihood Estimation algorithm(s)}}$



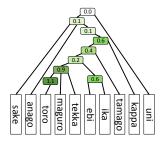
• Estimating τ given π_0 is tractable

Max Likelihood Estimation algorithm(s)



- ▶ Estimating τ given π_0 is tractable
 - ▶ by Dynamic Programming (DP) algorithm, similar to Matrix Chain Multiplication, Inside(-Outside) algorithm $\mathcal{O}(n^4)$
 - ightharpoonup contains θ_i estimation at each DP "partial solution"
- Estimating π_0 : Stochastic local search over π_0 space, similar to Simulated Annealing
 - 1. Sample ${\pi_0}^{\textit{new}}$ from proposal distribution current $P_{\tau,\theta}$
 - 2. Given π_0^{new} , find τ^{opt} , θ^{opt} by Dynamic Programming
 - 3. Bring to cannonical form $\Rightarrow \tau^{new}, \theta^{new} \succeq 0$
 - 4. Compute log-likelihood score, accept/reject like in Metropolis-Hastings, return to step 1

Experiments - Sushi preferences data



Data

N = 5000 permutations of n = 10 items Compared with:

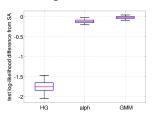
alph π_0 fixed, $\tau, \theta | \pi_0$ optimize

GM fixed τ , optimize π_0, θ

 ${
m HG}$ fixed au from [Huang,Guestrin,12], optimize heta

SA Simulated Annealing

Test set log-likelihood w.r.t SA



 $N_{test} = 300, N_{train} = 4700, 30$ replicates

"Sushi preference" data n = 12

types of sushi

"My top 3 preferences are ika, maguro, tekka, in this order" "I like uni least of all" "I prefer fish to non-fish"



. . .

Three good things about the RIM

- ▶ RIM is a general model (includes Mallows, generalized Mallows)
- ▶ likelihood $P(\pi|\tau(\vec{\theta}))$ factors according to tree (and partition function Z tractable)
- ▶ RIM has sufficient statistics

"Sushi preference" data n=12

types of sushi |ka| = |ka| = |ka| |ka| |ka| = |ka| |ka|

"Sushi preference" data n = 12

types of sushi

ika|maguro|tekka|{all other types} {all but ebi}|ebi

{sake,anago,...}|{tamago,ika,...}ှှု Ė₁

É2

Partial ranking σ [Huang & al, 10] $\sigma = (E_1|E_2|\dots|E_K)$ with

- $ightharpoonup E_1 \cup E_2 \cup \dots E_K = \operatorname{set}$ of items
- ightharpoonup shape $(n_1, \ldots n_K)$, $n_k = |E_k|, \sum n_k = n$

"Sushi preference" data n = 12

types of sushi

ika|maguro|tekka|{all other types}
{all but ebi}|ebi

$$\underbrace{\left\{ \text{sake}, \text{anago}, \dots \right\}}_{E_1} \mid \underbrace{\left\{ \text{tamago}, \text{ika}, \dots \right\}}_{E_2} \stackrel{\text{left}}{=} \underbrace{\left\{ \text$$

Partial ranking σ [Huang & al, 10] $\sigma = (E_1|E_2|...|E_K)$ with

- $E_1 \cup E_2 \cup \dots E_K = \text{set}$ of items
- ▶ shape $(n_1, ..., n_K)$, $n_k = |E_k|, \sum n_k = n$

Three good things about the RIM

- ▶ RIM is a general model (includes Mallows, generalized Mallows)
- ▶ likelihood $P(\pi|\tau(\vec{\theta}))$ factors according to tree ? YES [Huang et al, 10]
- ► RIM has sufficient statistics ? NO

Inferences with partial rankings in the RIM. Are they tractable?

The meaning of "tractable"

- **E**stimation of π_0 for RIM is intractable in the worst case
- ▶ We define tractable as $\mathcal{O}(N poly(n)) \times$ time (memory) for complete data

Inferences with partial rankings in the RIM. Are they tractable?

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Main technical difficulty

ightharpoonup marginal probability of a partial ranking σ

$$P(\sigma|\tau(\vec{ heta})) = \sum_{\pi \sim \sigma} P(\pi|\tau(\vec{ heta}))$$

where linear extension $\{\pi \sim \sigma\}$ of σ can have exponential size

Contributions

- 1. for marginal probability $P(\sigma|\tau(\vec{\theta}))$
 - exact formula and polynomial algorithm
 - proved algorithm no more than 2Nn more costly than for complete permutations (and sometimes much faster)
- 2. for pairwise marginals $E[Q_{ab}] = Pr[a \operatorname{precedes} b \mid \sigma, \tau(\vec{\theta})]$
 - exact recursive (polynomial) algorithm
 - proved algorithm no more costly than for complete permutations
- 3. for parameter $\vec{\theta}$ estimation (Maximum Likelihood)
 - ightharpoonup convex univariate minimization algorithm for each θi
 - ightharpoonup proved algorithm is $\mathcal{O}(Nn)$ more costly than for complete permutations
- 4. for structure search (Maximum Likelihood)

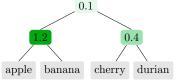
previous work

- complete data: local (simulated annealing) search algorithm with exact, tractable steps [Meek M 14]
- partial rankings: EM algorithm with approximate (or exponential) E step [Huang & al 10]

our contributions

- new "E step" based on completing the pairwise marginals $E[Q_{ab}]$
- algorithms above can use the completed pairwise marginals as if they were complete
 data

Computing the marginal probability $P(\sigma|\tau, \theta)$



$$P(a|b|c|d) \propto e^0$$

 $P(b|a|c|d) \propto e^{-1.2}$
 $P(c|b|a|d) \propto e^{-1.2-2\times0.1}$

RIM probability for complete data $P(\pi|\tau, \vec{\theta})$ (with v_i = number of inversions of π_0 at node i)

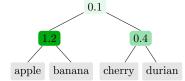
$$P_{ au, ec{ heta}}(\pi) = \prod_{i \in nodes} rac{\mathrm{e}^{- heta_i \mathbf{v}_i}}{G_{L_i, R_i}(\mathrm{exp}(- heta_i))}$$

with
$$G_{L,R}(q) = \frac{(q)_{L+R}}{(q)_L(q)_R}$$
, $(q)_n = \prod_{i=1}^n (1-q^i)$.

RIM probability for partial ranking σ [M, Meek in prep]

$$P_{\tau,\vec{\theta}}(\sigma) = \prod_{i \in nodes} (factor at node i)$$

Marginal $P(\pi| au, ec{ heta})$ for partial ranking σ



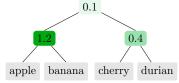
Sufficient to consider root node Complete ranking $\pi = (c|a|b|d)$

factor =
$$\frac{e^{-2\theta}}{G_{2,2}(e^{-\theta})}$$

Partial ranking
$$\sigma = (c | \{a, b, d\})$$

$$\mathsf{factor} = \frac{e^{-2\theta} \, {\color{red}\mathsf{G}_{0,1}(e^{-\theta}) \, {\color{red}\mathsf{G}_{2,1}(e^{-\theta})}}}{{\color{red}\mathsf{G}_{2,2}(e^{-\theta})}}$$

Marginal $P(\pi| au, \vec{ heta})$ for partial ranking σ



Sufficient to consider root node Complete ranking $\pi = (c|a|b|d)$

Partial ranking
$$\sigma = (c | \{a, b, d\})$$

factor =
$$\frac{e^{-2\theta}}{G_{2,2}(e^{-\theta})}$$

factor =
$$\frac{e^{-2\theta} G_{0,1}(e^{-\theta}) G_{2,1}(e^{-\theta})}{G_{2,2}(e^{-\theta})}$$

In general, at some internal node where

- ightharpoonup set $\mathcal L$ is merged with set $\mathcal R$
- ▶ partial ranking σ restricted to $\mathcal{L} \cup \mathcal{R}$ is $E_1|E_2|\dots|E_K$ with $E_k = L_k \cup R_k$, $L_k \subseteq \mathcal{L}$, $r_k \subseteq \mathcal{R}$
- factor of $P(\sigma|\tau(\vec{\theta}))$ at this node is

$$g(I_{1:K}, r_{1:K}, \theta) = \frac{e^{-\theta v} G_{I_1, r_1}(e^{-\theta}) G_{I_2, r_2}(e^{-\theta}) \dots G_{I_K, r_K}(e^{-\theta})}{G_{|\mathcal{L}|, |\mathcal{R}|}(e^{-\theta})}$$

where ${\it v}=\#$ inversions in σ at node $\leq\#$ inversions in $\pi\sim\sigma$

Marginal $P(\pi|\tau, \vec{\theta})$ – how much extra computation?

How many additional factors?

Marginal $P(\pi|\tau,\theta)$ How many additional Rem 1 $G_{0,r}=G_{I,0}=1$ Rem 2 at each node, a Marginal $P(\pi|\tau, \vec{\theta})$ – how much extra computation?

How many additional factors?

Rem 2 at each node, at least one of L_k , R_k decreases (and their initial sum is n)

▶ Hence, no more than n-1 extra factors (but sometimes much fewer)

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How many additional factors?

Rem 2 at each node, at least one of L_k , R_k decreases (and their initial sum is n)

- ▶ Hence, no more than n-1 extra factors (but sometimes much fewer)
- ► Example top-t rankings $\sigma = (ika|maguro|sake|{everything else}) P(\sigma|\tau, \vec{\theta})$ has at most t-1 non-trivial factors

How many additional factors?

Rem 1 $G_{0,r} = G_{l,0} = 1$

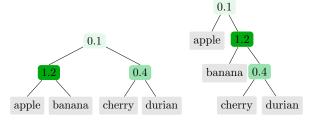
Rem 2 at each node, at least one of L_k , R_k decreases (and their initial sum is n)

- ▶ Hence, no more than n-1 extra factors (but sometimes much fewer)
- Example top-t rankings $\sigma = (ika|maguro|sake|\{everything else\}) P(\sigma|\tau, \vec{\theta})$ has at most t-1 non-trivial factors

How much additional computation?

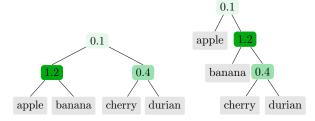
- $G_{L,R}$ is computed recursively over I = 0, ... L, r = 1, ... R
- lacktriangle Hence, all $G_{l,r}(heta)$ in numerator are cached while computing the denominator
- \blacktriangleright Overhead for whole sample of size N is no more than nN lookups+multiplications
- ▶ For comparison, for a complete whole sample
 - ▶ computation of sufficient statistics is $\mathcal{O}(n^2N)$
 - ▶ computation of Z given $\vec{\theta}$ is $\mathcal{O}(n^2 \log n)$

Independence properties



- define $Q_{ab} = 1$ iff a precedes b
- $lackbox{ }Q_{ab}\perp Q_{cd}$ whenever $\mathsf{path}(a,b)\cap \mathsf{path}(c,d)=\emptyset$

Independence properties



- define $Q_{ab} = 1$ iff a precedes b
- ▶ $Q_{ab} \perp Q_{cd}$ whenever $\mathsf{path}(a,b) \cap \mathsf{path}(c,d) = \emptyset$
- Indepence checking can reveal the "branching structure" (but not π_0)
- \blacktriangleright In progress: combine independence tests with local search to estimate τ

Conclusion: No need to compromise!

Goals of inference in models on permutations

- ► Flexible w.r.t observation model (i.e. input data)
 - partial rankings, pairwise observations
- ► Flexible w.r.t generative model
 - ▶ RIMs are a class of flexible, identifyable, intepretable models
- ▶ Exact and tractable algorithms, closed form expression