

BIOSTAT 524

4/17/23

Lecture 7

NP and Hierarchical
clustering

(lecture given on UW tablet)

L IV.1 - Hierarchical
LV SVM

Lecture IV – Non-parametric clustering

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*Think of next topics
for rest of 527*

- LIV.1 Hierarchical
- LV SVM Kernel
- HW2 Wednesday

- 1 Paradigms for clustering ✓

- 2 Methods based on non-parametric density estimation ✓

cluster tree (level sets)
DB scan

- 3 [Model-based: Dirichlet process mixture models]

Reading AoNPS Ch.: –, HTF Ch.: 14.3 Murphy Ch.: 11.[1], 11.2.1-3, 11.3, Ch 25

The Nugent-Stuetzle algorithm

Algorithm Nugent-Stuetzle

Input Data $\mathcal{D} = \{x_i\}_{i=1:n}$, kernel $K(z)$

① Compute KDE $f(x)$ for chosen h

② for levels $0 < l_1 < l_2 < \dots < l_r < \dots < l_R \geq \sup_x f(x)$

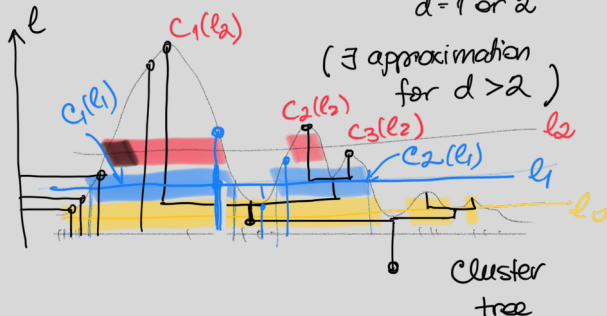
① find level set $L_r = \{x \mid f(x) \geq l_r\}$ of f

② if L_r **disconnected** then each connected component is a cluster $\rightarrow (C_{r,1}, C_{r,2}, \dots, C_{r,K_r})$

Output clusters $\{(C_{r,1}, C_{r,2}, \dots, C_{r,K_r})\}_{r=1:R}$

$f(x^i)$ $i=1:n$
sort by $f(x^i)$

tractable only in
 $d=1$ or 2

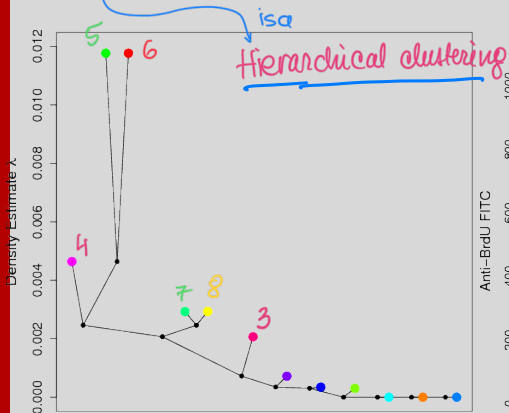


Remarks

- every cluster $C_{r,k} \subseteq$ some cluster $C_{r-1,k'}$
- therefore output is hierarchical clustering
- some levels can be pruned (if no change, i.e. $K_r = K_{r-1}$)
- algorithm can be made recursive, i.e. efficient
- finding level sets of f tractable only for $d = 1, 2$
- for larger d , $L_r = \{x_i \in \mathcal{D} \mid f(x_i) \geq l_r\}$
- to find connected components
 - for $i \neq j \in L_r$
 if $f(tx_i + (1-t)x_j) \geq l_r$ for $t \in [0, 1]$
 then $k(i) = k(j)$
- confidence intervals possible by resampling

} heuristic
for $d > 2$

Cluster tree with 13 leaves (8 clusters, 5 artifacts)

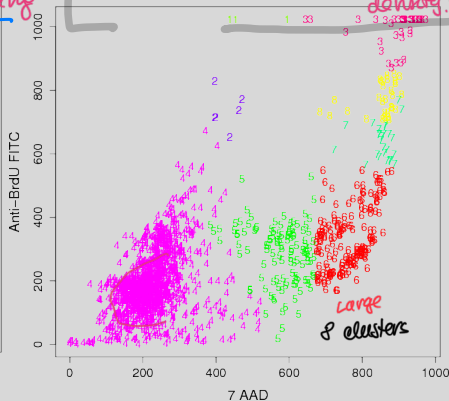


(from ?)

Rebecca
Nugent

Parenthesis on plotting

- default markers → to see outliers!
- really small dots → to see density!



Chaudhuri-Dasgupta Algorithm

- Uses k -nearest neighbor graphs (filtration)
- Parameters k (nearest neighbors) and $\alpha \in [1, 2]$
- for $r \geq 0$, $G_r = (V_r, E_r)$ with
 - $x_i \in V_r$ iff distance to k -nn of $x_i \leq r$
 - $(x_i, x_j) \in E_r$ iff $\|x_i - x_j\| \leq \alpha r$

Consistency Theorem For any ϵ (separation parameter) and δ (confidence), $\alpha \in [\sqrt{2}, 2]$ (graph density), if $k = C \log^2(1/\delta) \frac{d \log n}{\epsilon^2}$ for any two clusters C, C' in cluster tree, there exists a level r so that $C \cap D, C' \cap D$ are clusters at level r

↑
note $k \sim d \log n$
Not small

NP anything
 k -NN must have
large enough k

↑
GENERAL
RULE OF THUMB
FOR NP

The K-nn density estimator

The K-nn density estimator

- Let $B_r(x)$ be the (closed) ball of radius r centered at x
- If $|B_r(x^i) \cap \mathcal{D}| = k$ then $\hat{p}(x^i) = \frac{1}{r^n \omega_n} \frac{k}{n}$ is an estimate of the density at x^i
 - $\omega_n = \pi^{n/2} / \Gamma(n/2 + 1)$ is the volume of the unit ball in $\mathbb{R}^n$
 - intuitively, the ball of radius r contains k/n probability mass
 - Note that the density $\hat{p}$ is not required to integrate to 1



DBScan (Heuristic)

- Introduced with no proof, but widely used. Implicitly based on the K-nn estimator
- Parameters r radius, m minimum number points
- Definitions **core** $Q = \{x^i \in \mathcal{D}, \text{ with } |B_r(x^i) \cap \mathcal{D}| \geq m\}$
- border** $B = \{x^i \in \mathcal{D} \setminus Q, \text{ so that } x^i \in B_r(x^j), x^j \in Q\}$
- outliers (noise)** $O = \mathcal{D} \setminus (Q \cup B)$



$r = \text{neighborhood radius}$
 $m = 4$

$m \uparrow$ fewer core
 more clusters?
 ?

- Algorithm idea
- Construct directed graph G with edges (i, j) where $x^i \in Q, j \in B_r(x^i)$ $r \uparrow \Rightarrow$ same $m \Rightarrow$ more core
- The graph edges between core points are undirected/symmetric, the other are from core to border $\Rightarrow$ efficient!
- Clusters are determined by the connected components of the graph restricted to Q . fewer clusters?
- The border points are assigned to a cluster containing x^i so that $x^i \in B_r(x^i), x^j \in Q$ Note that this assignment is not unique!
- Heuristic algorithm estimates r, m

[Supplement: Chaudhuri-Dasgupta Algorithm]

Consistency Theorem For any ϵ (separation parameter) and δ (confidence), $\alpha \in [\sqrt{2}, 2]$ (graph density), if $k = C \log^2(1/\delta) \frac{d \log n}{\epsilon^2}$

for any two clusters C, C' in cluster tree, there exists a level r so that $C \cap \mathcal{D}, C' \cap \mathcal{D}$ are clusters at level r

- r depends on λ = "bridge" between C, C' (and $\sigma > 0$ "tube" width)

$$r^d \omega_d \lambda = \frac{k}{n} + \dots \text{confidence term}$$

- it follows that the needed sample size n at level λ

$$n = \mathcal{O} \left(\frac{d}{\lambda \epsilon^2 (\sigma/2)^d \omega_d} \log \frac{d}{\lambda \epsilon^2 (\sigma/2)^d \omega_d} \right)$$

- this **sample complexity** n is almost tight
- for $\alpha < \sqrt{2}$ sample complexity is exponential in d
- New results [Kent, B. P., Rinaldo, A. and Verstynen, T. 2013]
- **Remark:** algorithm(s) can be applied in **any metric space**

Lecture IV – Hierarchical clustering. (Comparing clusterings)

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Hierarchical Methods of Clustering

(Herovec Data mining)

Agglomerative (bottom up):

- Initially, each point is a cluster
- Repeatedly combine the two "nearest" clusters into one

n^3

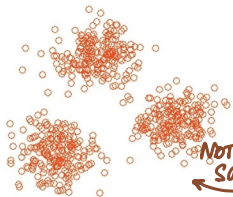
greedy algorithm

Divisive (top down):

- Start with one cluster and recursively split it

Dendrogram

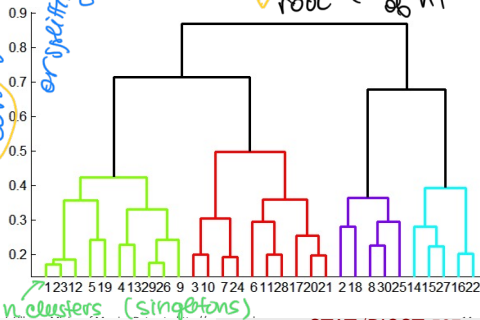
root (1 cluster of n points)



NOT
same
data

"leaves"

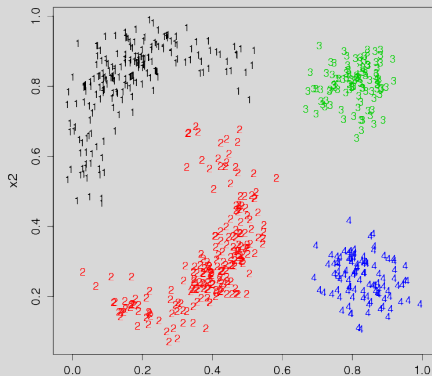
cost of merging
or splitting



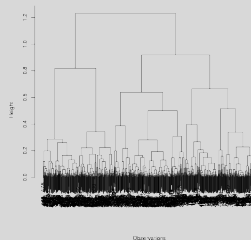
NOT ROBUST: single
— distance btw points $\Rightarrow$ linkage
* Cost e.g. $\sum \|x_i - x_j\|^2$ means
not in same cluster

What is hierarchical clustering?

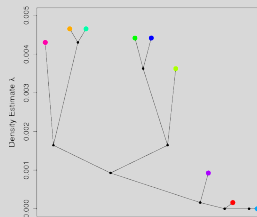
- Clusters have cluster structure
- Represented by
 - **Dendrogram**
 - **Cluster Tree** (only from KDE)



Dendrogram



Cluster Tree

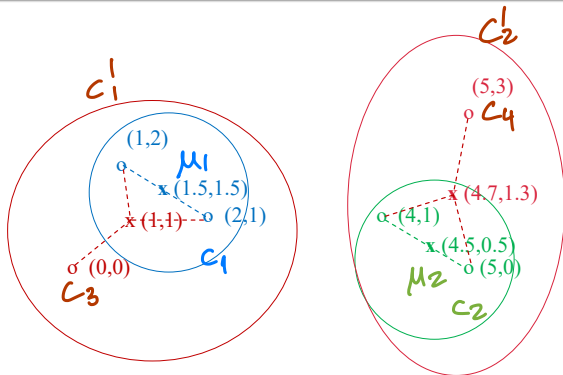


Hierarchical clustering – Overview

(Dendrograms)

- **Agglomerative** (bottom up)
 - **Single linkage**
 - based on Minimum Spanning Tree
 - $\mathcal{O}(n^2 \log n)$
 - sensitive to outliers
 - Heuristics – average linkage
 - **Agglomerative K-means**
 - Loss $\mathcal{L}(\Delta_K) = 0$ for $K = n$
 - When $K \leftarrow K - 1$ (two clusters merged), $\mathcal{L}$ increases
 - For $K = n, n - 1, \dots, 2$, iteratively merge the 2 clusters that minimize increase of $\mathcal{L}$
 - $\mathcal{O}(n^3)$ – too expensive for big data
- **Divisive** (bottom down)
 - Recursively split $\mathcal{D}$ into $K = 2$ clusters
 - almost any clustering algorithm (e.g. K-means, min diameter)
 - notable example **Spectral clustering** (later)
 - Advantages
 - most important splits are first
 - can stop after only a few splits

Example: Hierarchical clustering



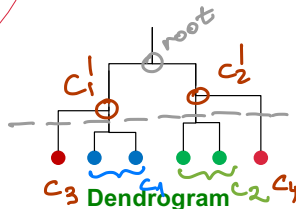
$$C_1' = C_1 \cup C_3$$

$$C_2' = C_2 \cup C_4$$

$$\text{root} = C_1' \cup C_2'$$

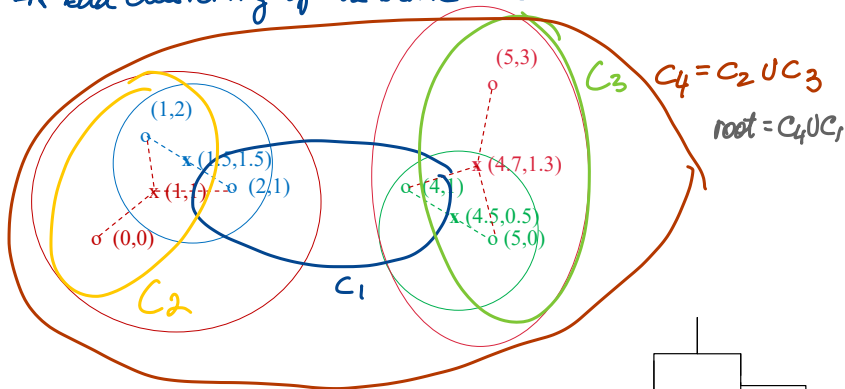
Data:

o ... data point
x ... centroid



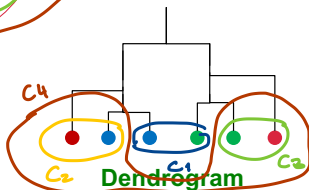
Example: Hierarchical clustering

A bad clustering of the same data

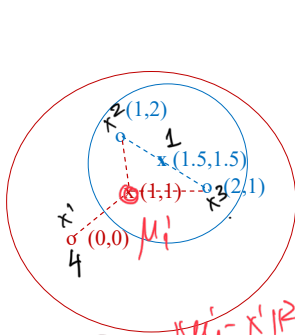


Data:

- o ... data point
- x ... centroid



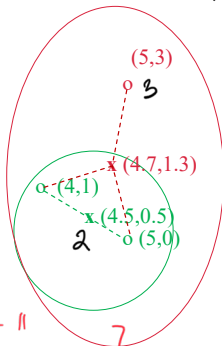
Example: Hierarchical clustering



$$\frac{1}{\sigma^2 2\pi} \left[e^{-\frac{\| \mu^1 - x^1 \|^2}{2\sigma^2}} + e^{-\frac{\| \mu^1 - x^2 \|^2}{2\sigma^2}} + e^{-\frac{\| \mu^1 - x^3 \|^2}{2\sigma^2}} + e^{-\frac{\| \mu^1 - x^4 \|^2}{2\sigma^2}} \right] + \dots$$

Data:

o ... data point
x ... centroid



Mix with k clusters

$k=4$

$$\text{likelihood} = \frac{1}{\sigma^2 2\pi} e^{-\frac{\| \mu^1 - x^1 \|^2}{2\sigma^2}} + \dots$$

$$+ \frac{1}{\sigma^2 2\pi} \left[e^{-\frac{\| \mu^1 - x^2 \|^2}{2\sigma^2}} + e^{-\frac{\| \mu^1 - x^3 \|^2}{2\sigma^2}} \right]$$

